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(54) Title of Invention

Data processing system

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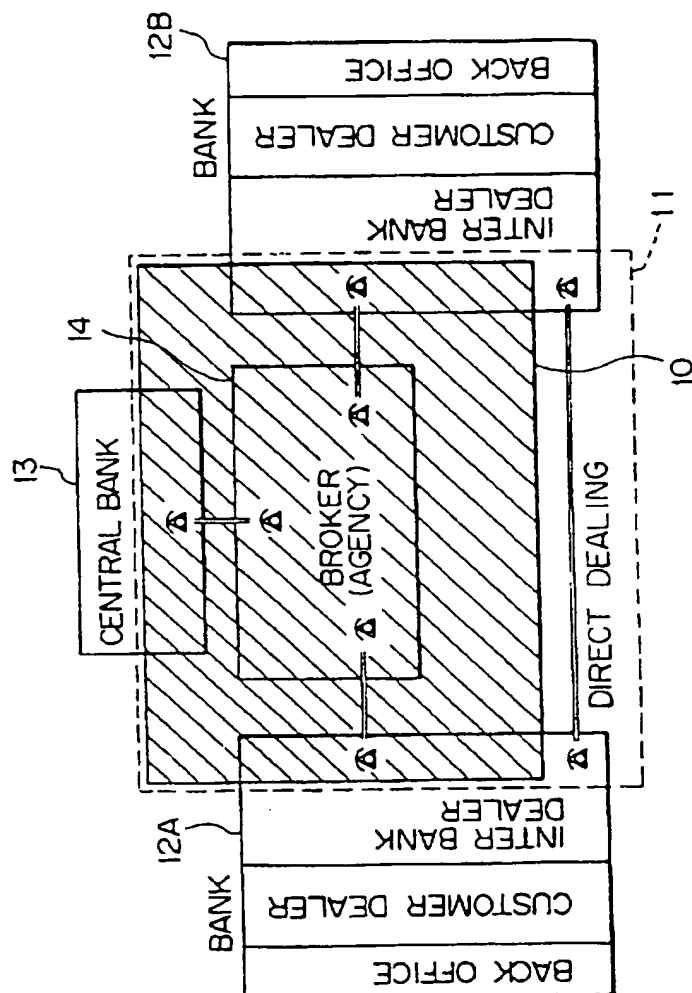
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(22) Date of filing 17.03.1994	(73) Proprietor(s) Fujitsu Limited  (Incorporated in Japan)
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(56) Documents cited None	Nanwa-Nihonbashi Building 4-2-16 Nihonbashi Muromachi Chuo-Ku Tokyo 103 Japan
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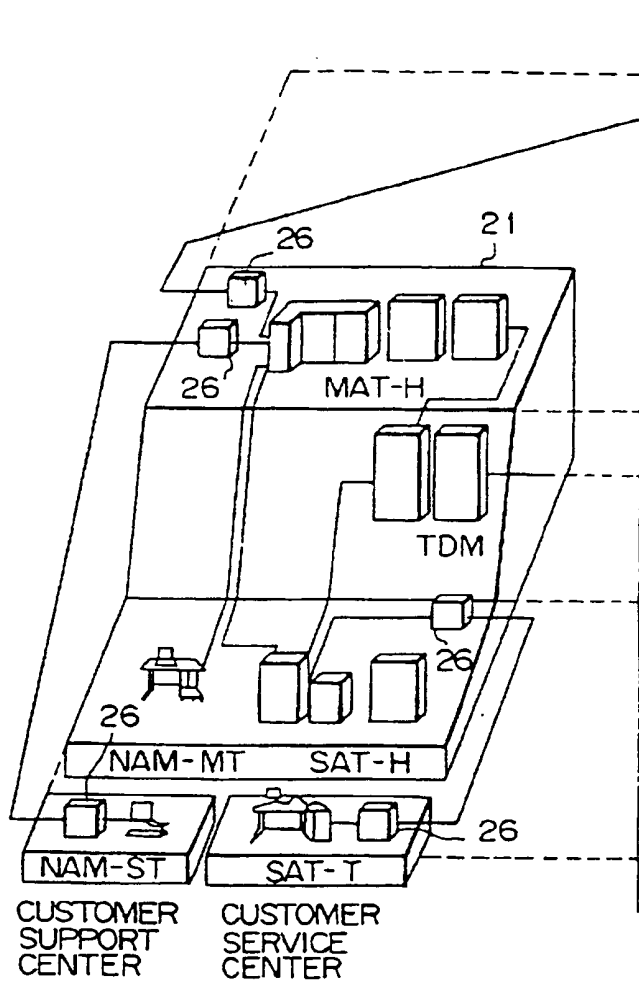
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Fig. 1



$\frac{2}{26}$ *Fig.2**Fig.3'*

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Fig.3

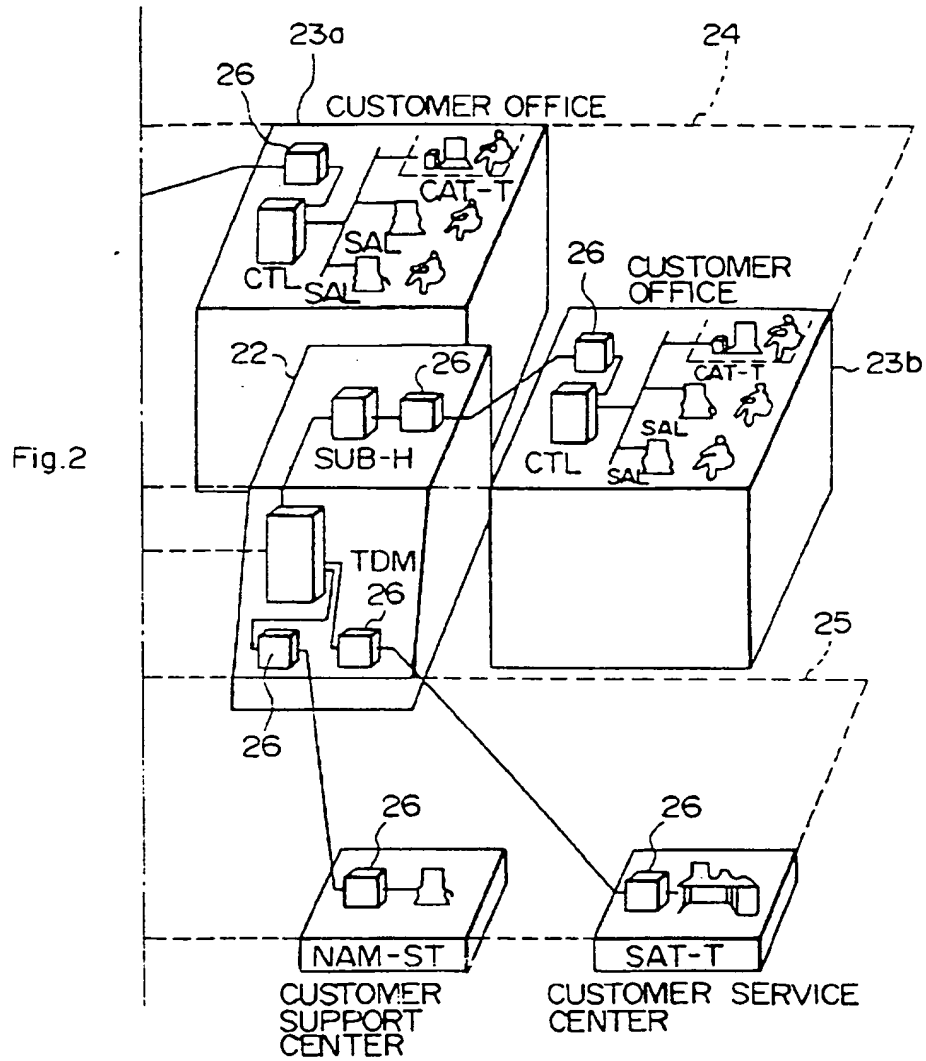
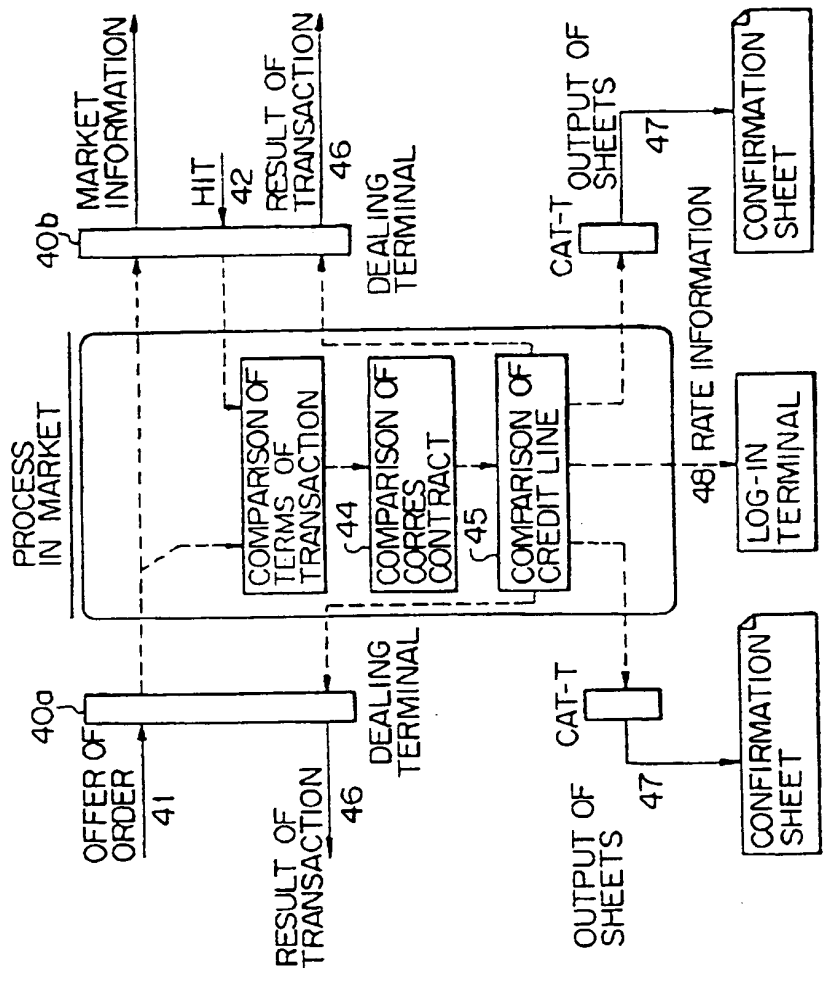


Fig.4



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Fig.5

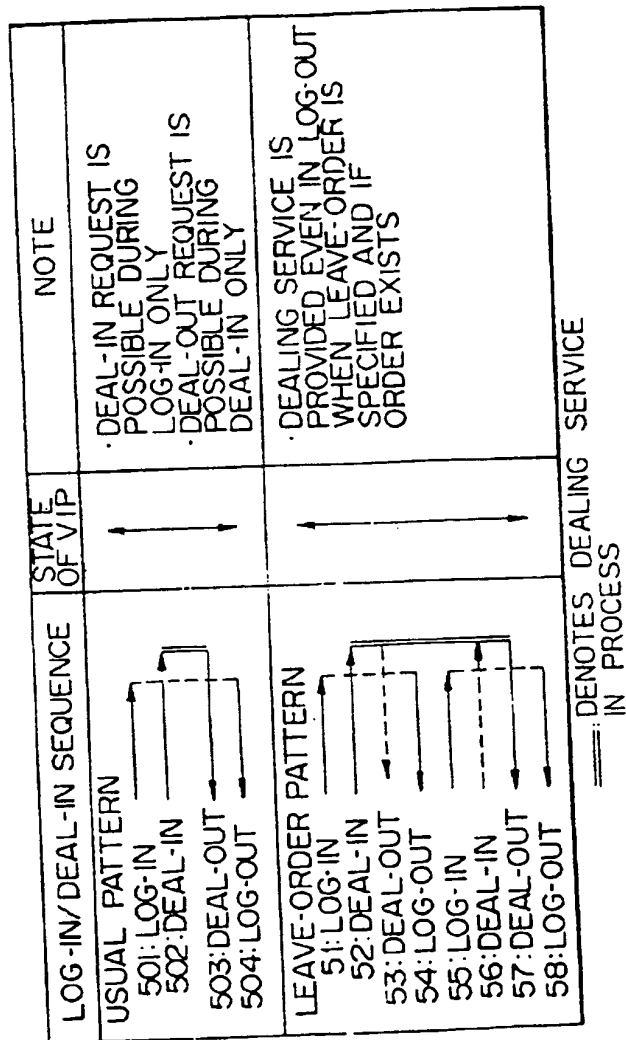


Fig.6

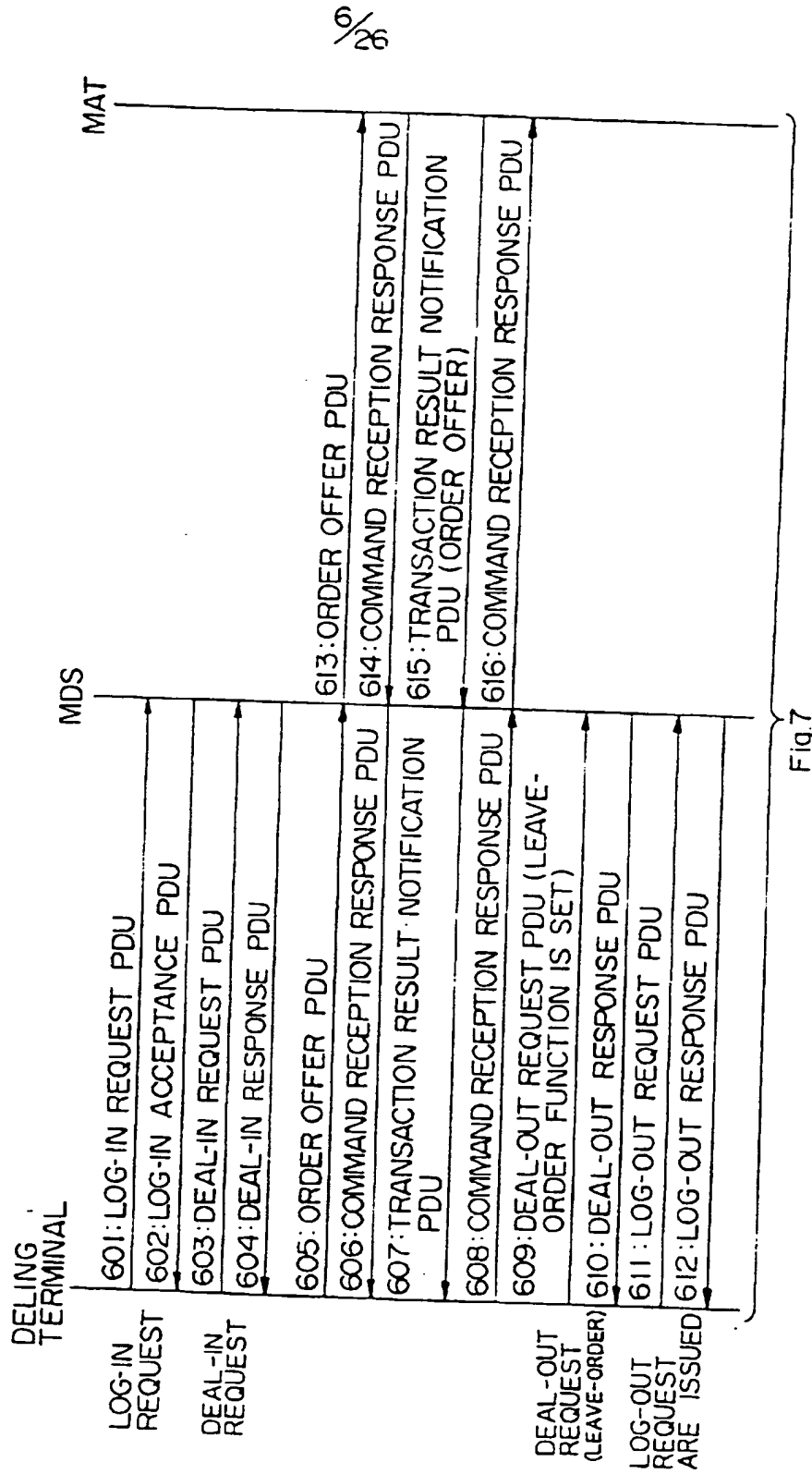
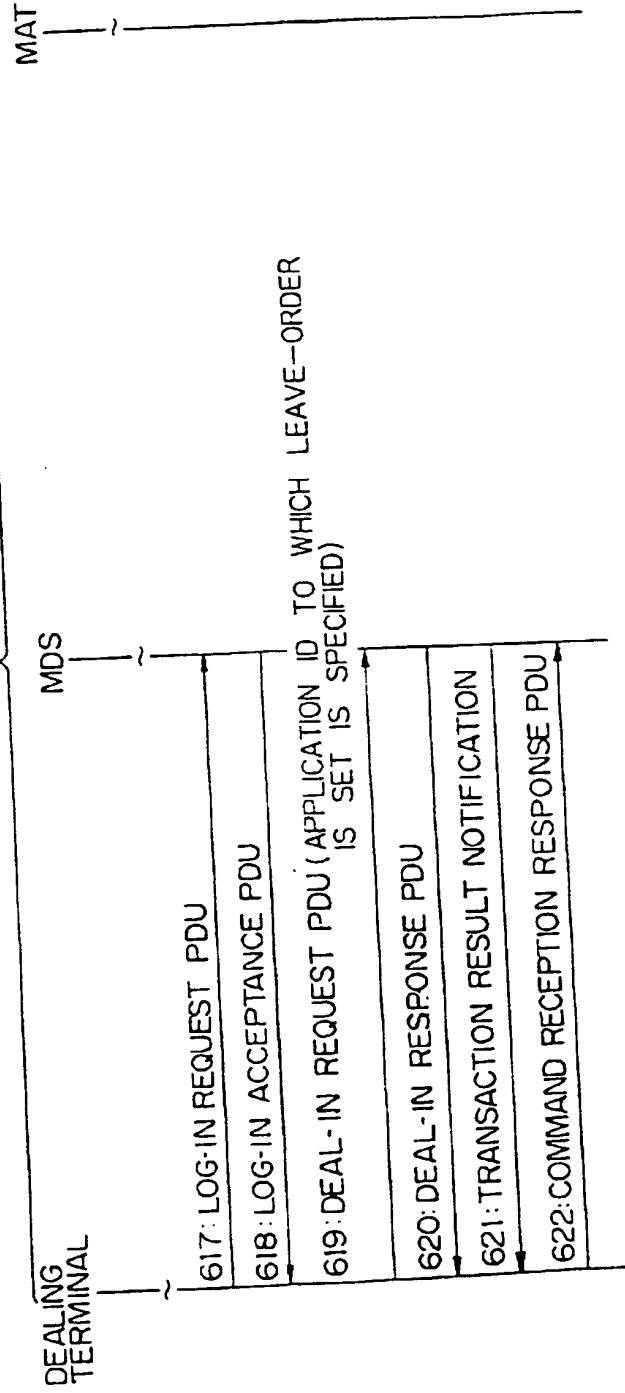


Fig.7

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Fig.7

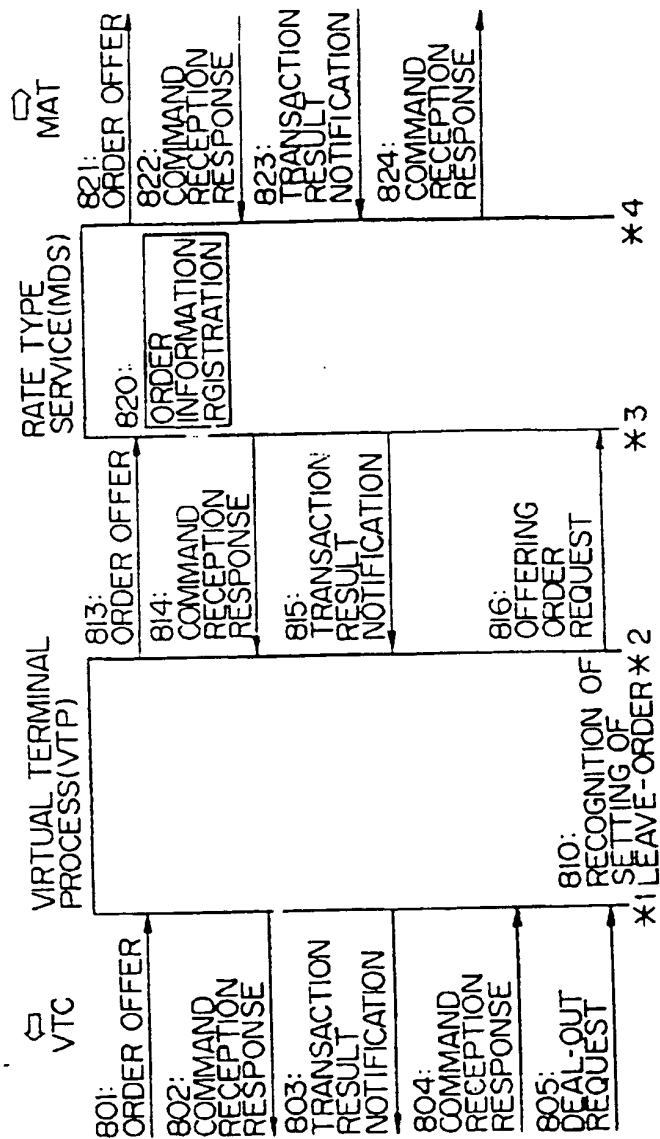
Fig.6





$\frac{a}{26}$ 

Fig.8A



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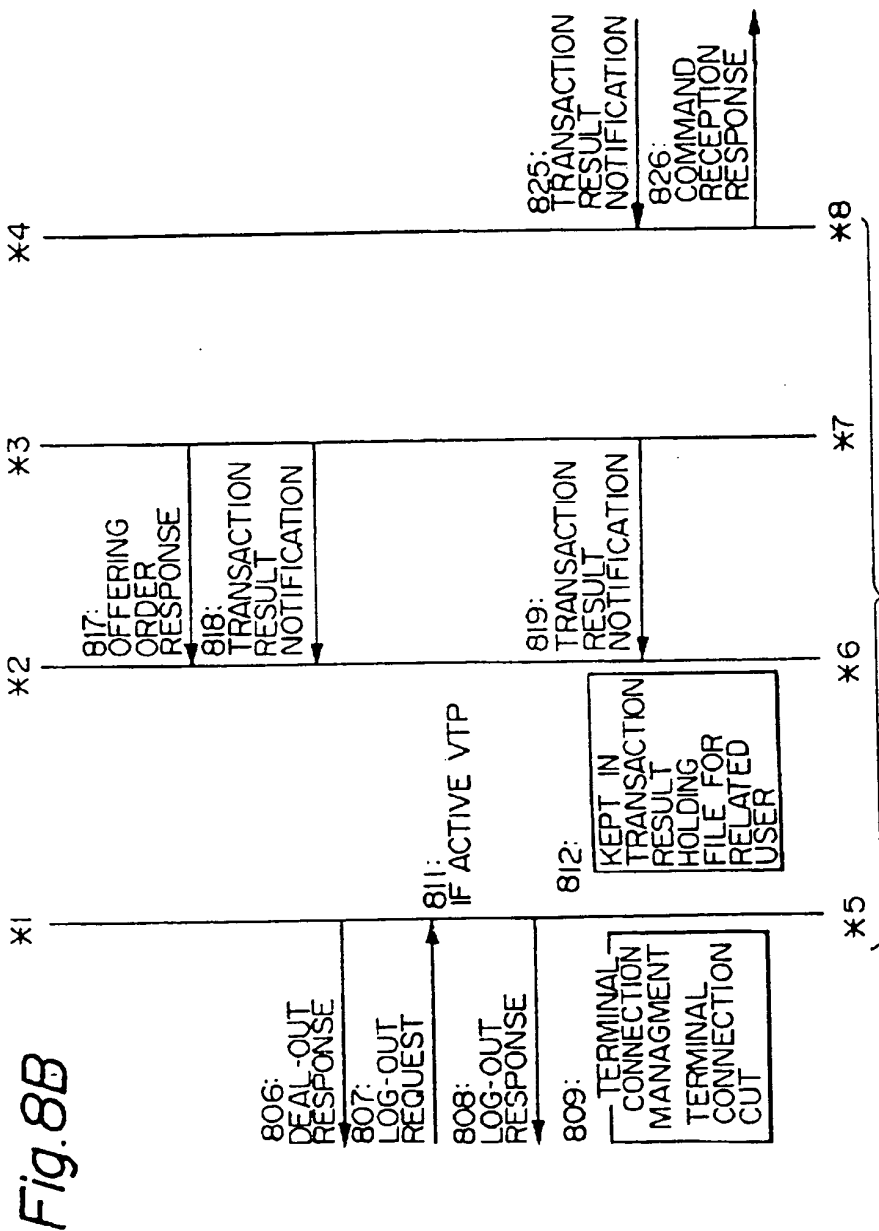
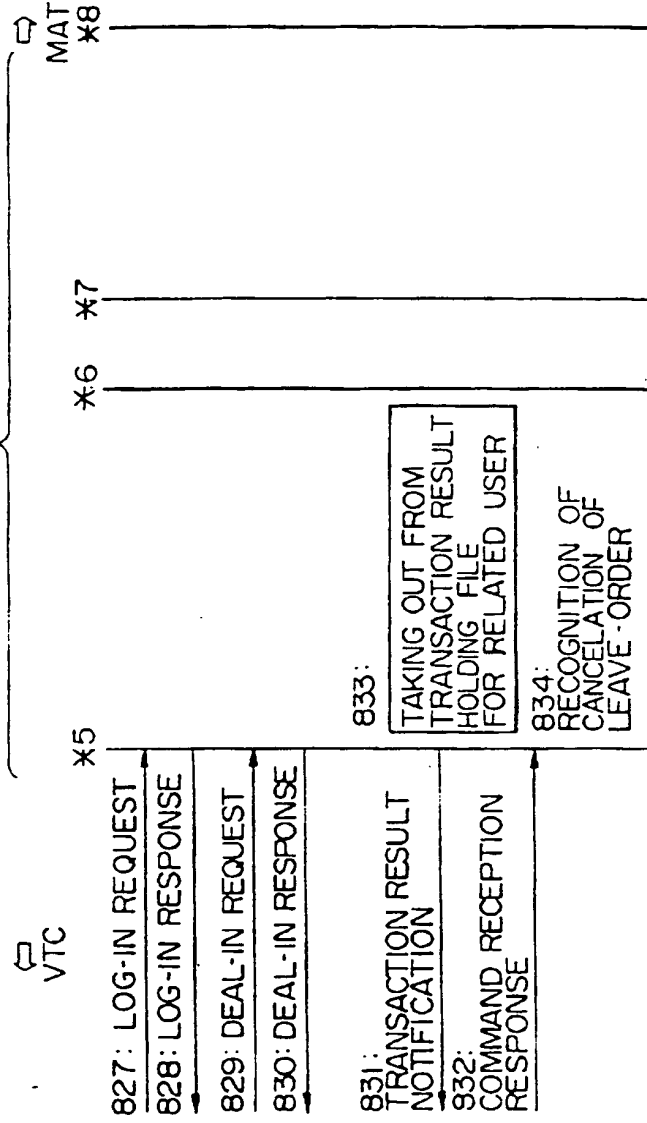


Fig. 9

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Fig.9

Fig.8B





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Fig.10B

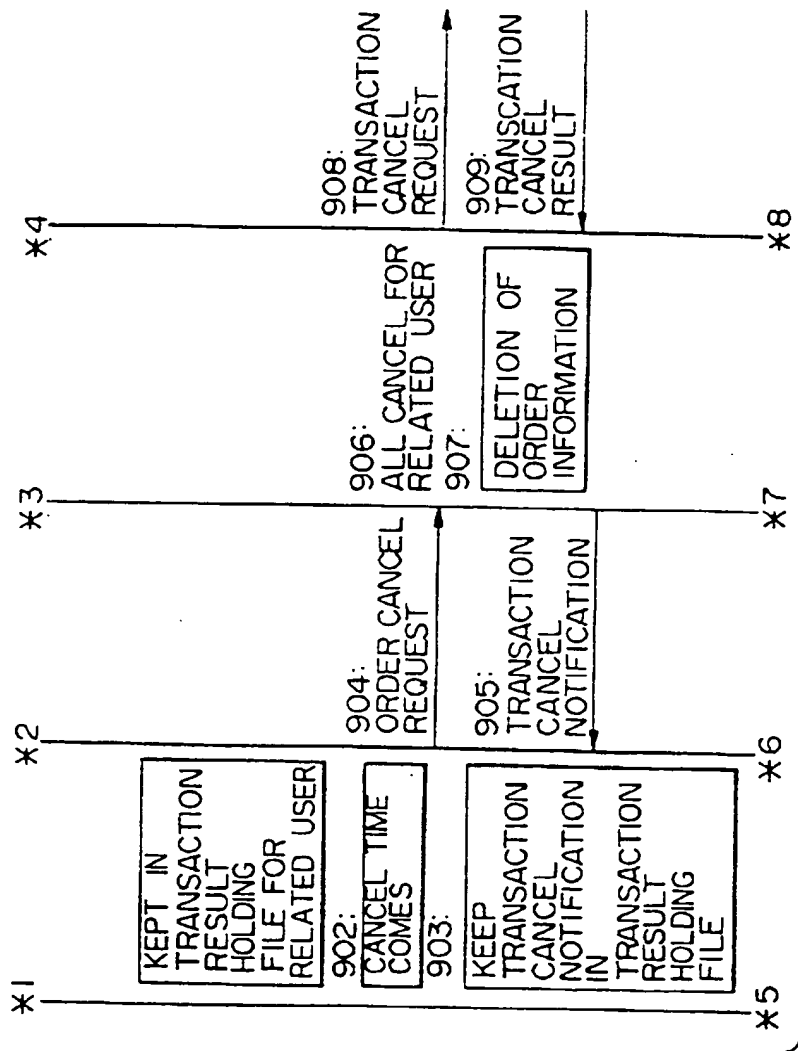
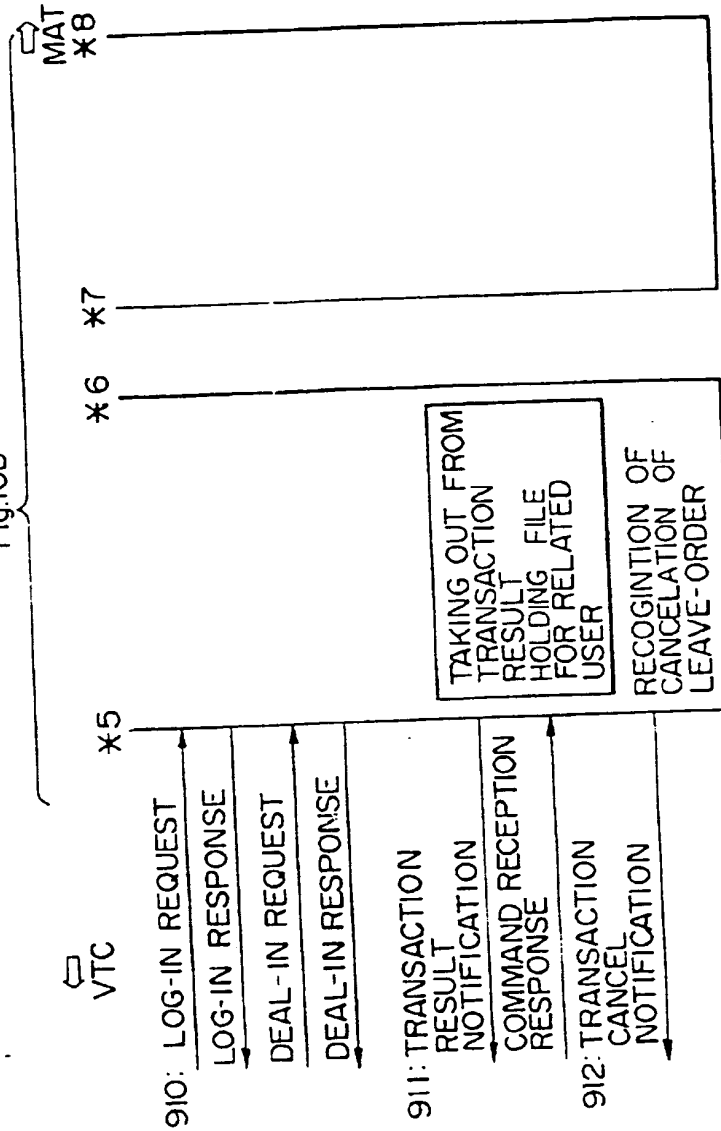


Fig.11

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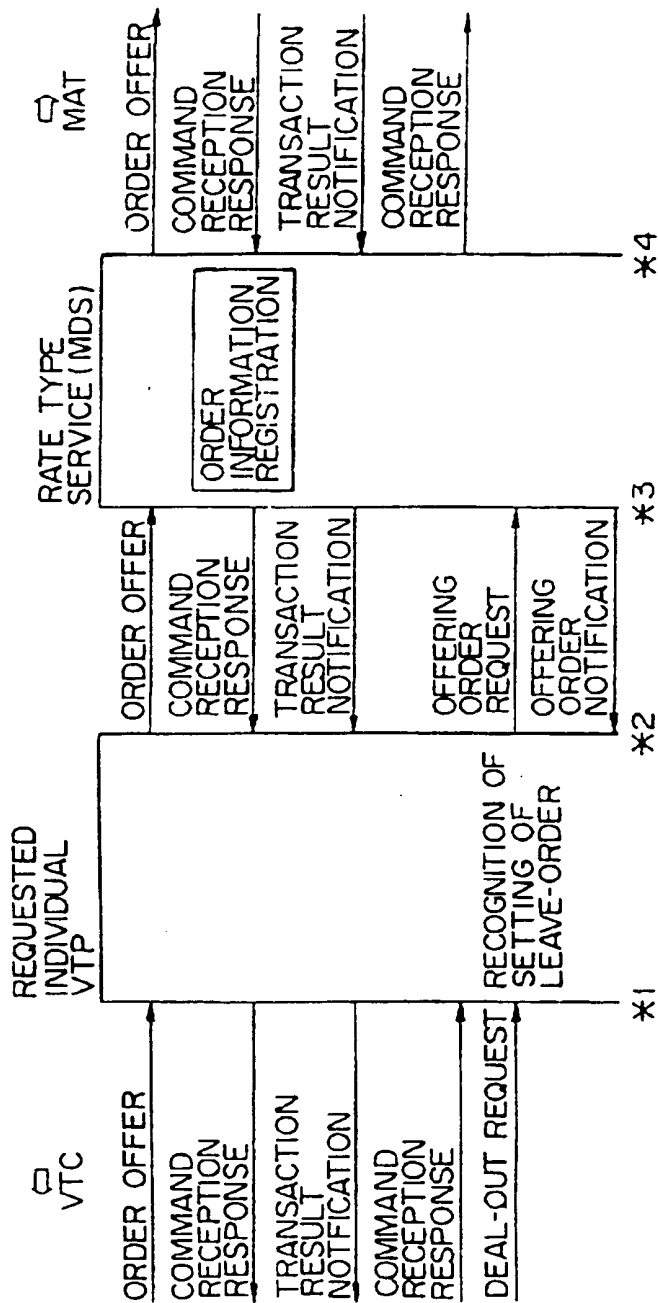
Fig.11

Fig.10B



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Fig.12A



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Fig.12B

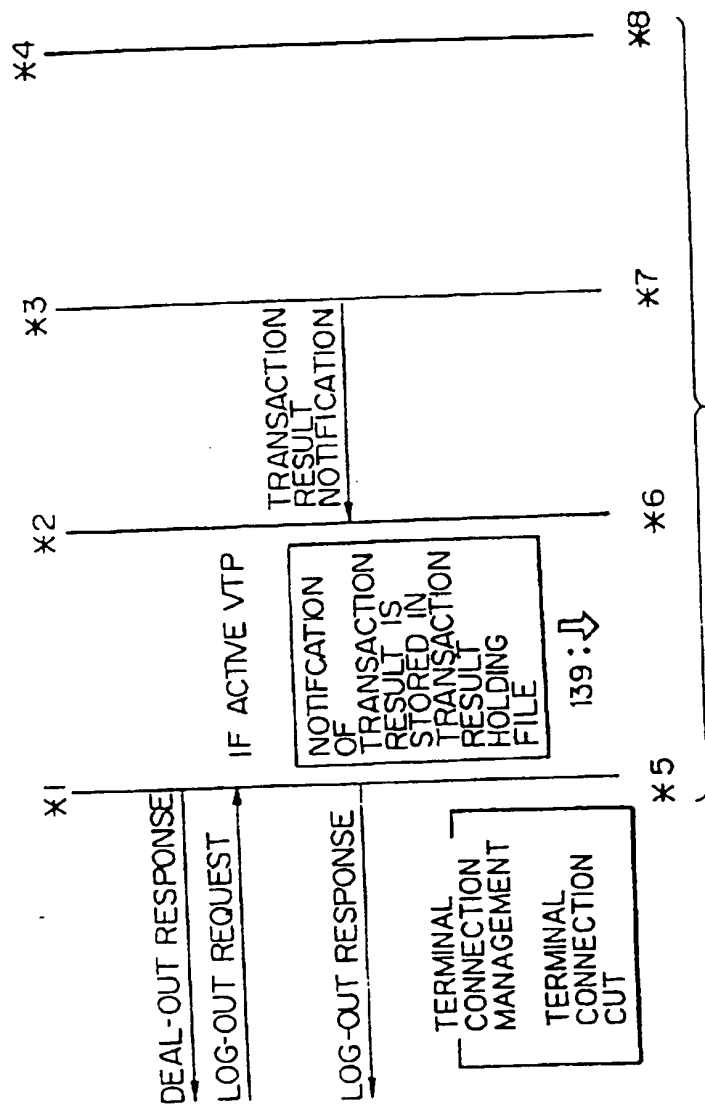


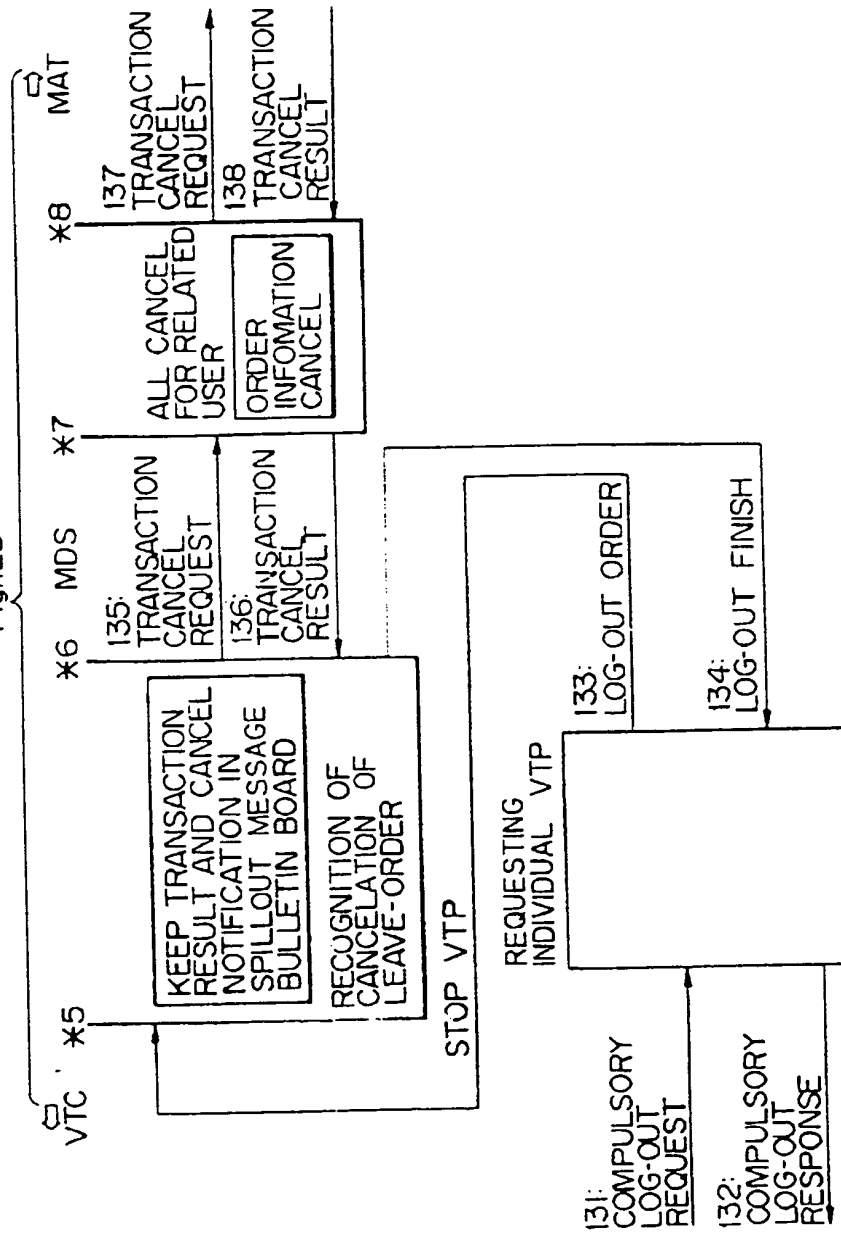
Fig.13

TERMINAL  
CONNECTION  
MANAGEMENT  
TERMINAL  
CONNECTION  
CUT



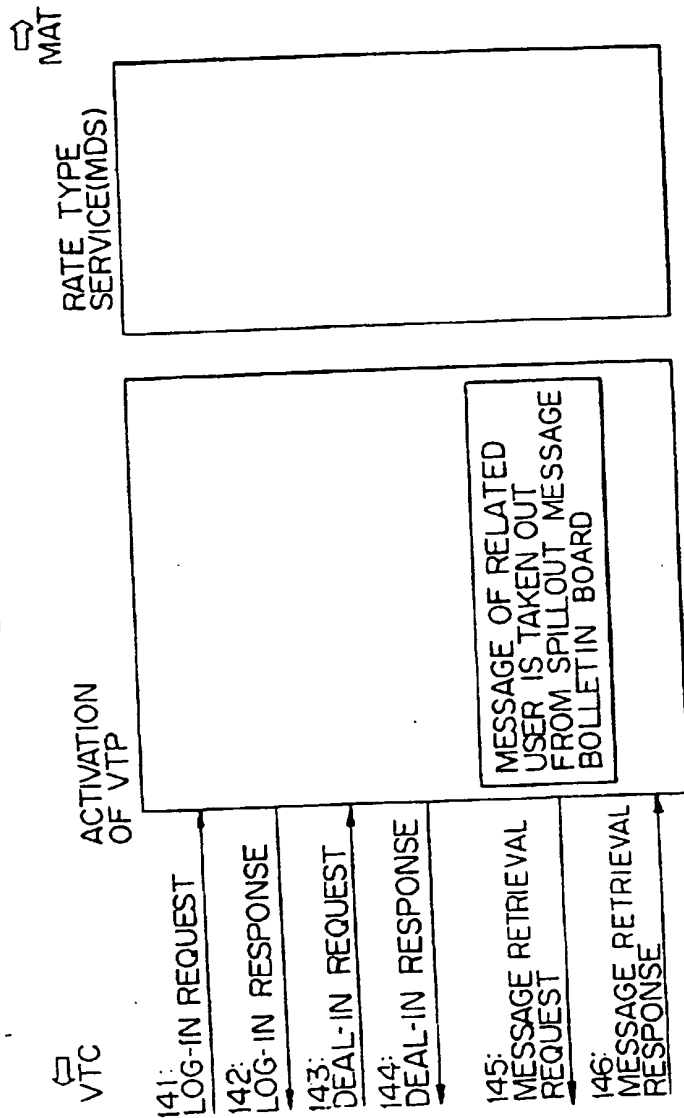
Fig.13

Fig.12B



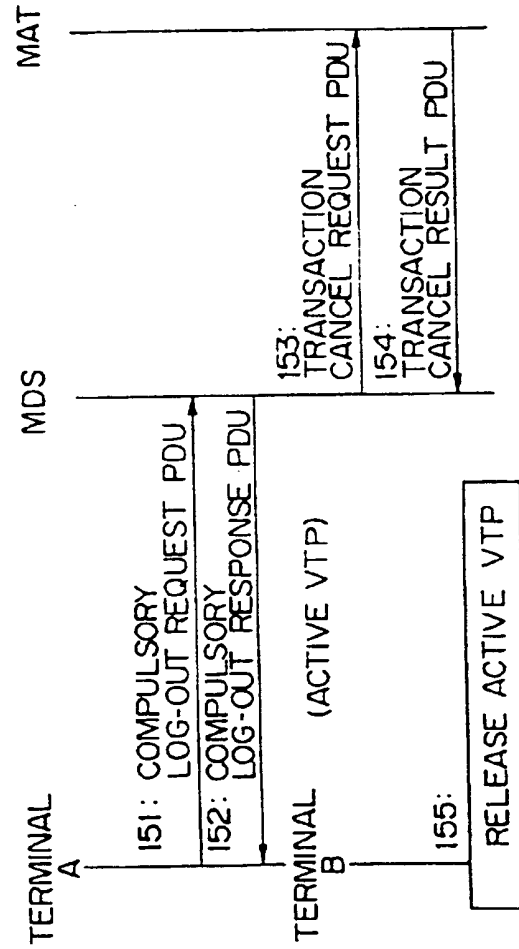
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Fig.14



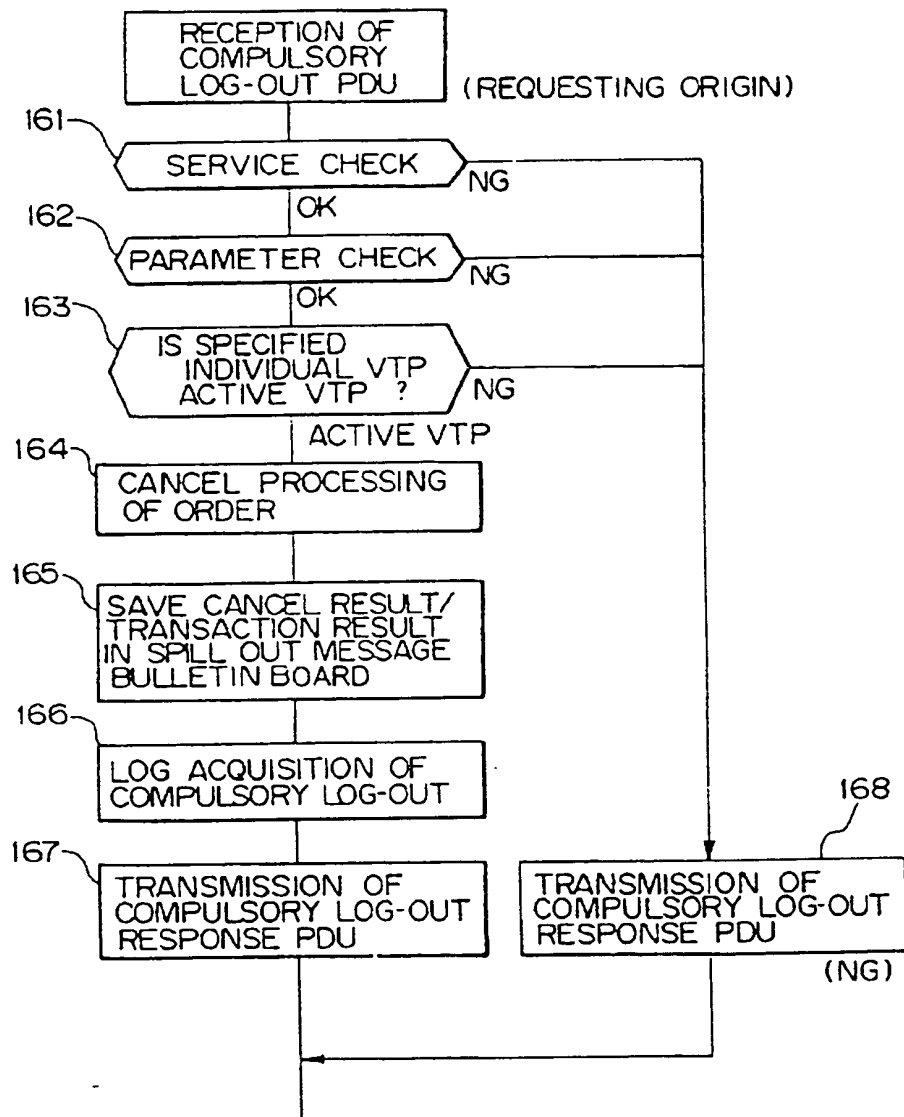
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Fig.15



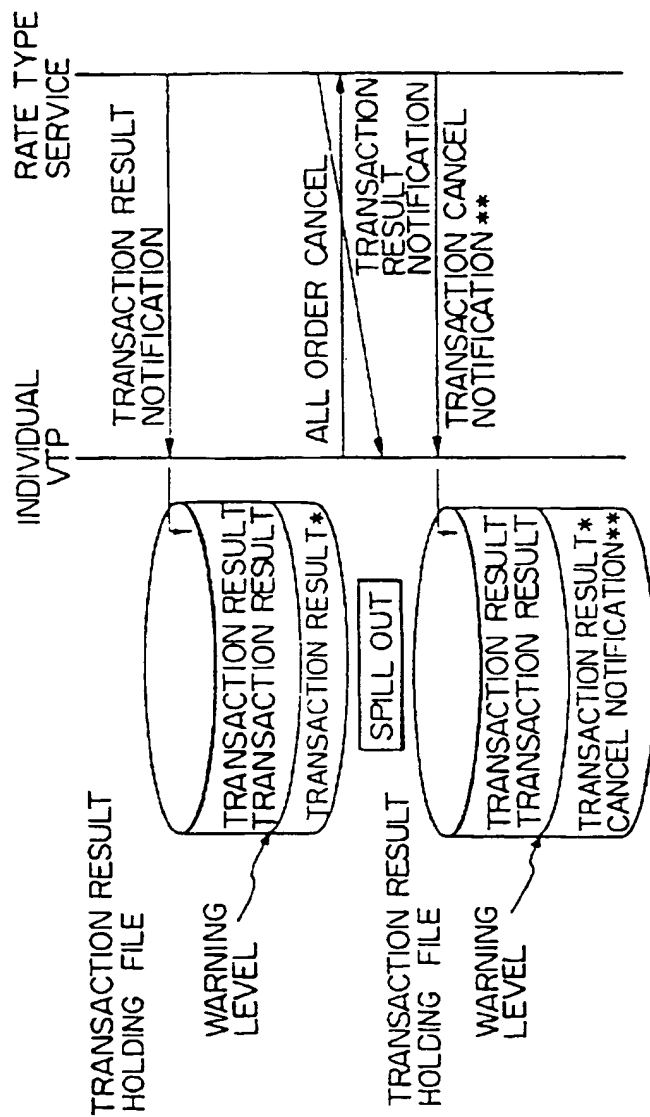
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Fig.16



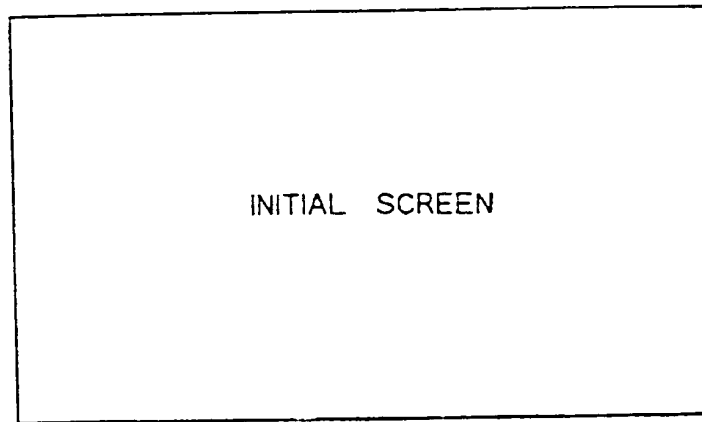
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Fig.17



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26

*Fig.18A*



*Fig.18B*

—	SERVICE LOG-IN	
<input type="checkbox"/>	USER ID	<input type="text"/>
<input type="checkbox"/>	PASSWORD	<input type="text"/>
<input type="text"/>		EXEC QUIT

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Fig.19A

MANAGEMENT	
<input type="checkbox"/>	DEALING SERVICE
<input type="checkbox"/>	CONFIRMATION SERVICE
<input type="checkbox"/>	ALL MARKET INFORMATION

MENU	1993/04/12 10:00 M	<input type="checkbox"/>					
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Fig.19B

DEALING SERVICE		/VAL 1993/04/14	
MARKET AMT	PTY	YOUR OFFERS	
		/ *	
			BUY/SELL
			TOTAL AMT
			TOTAL TRS
			AVE. PRICE

MENU	1993/04/12 10:02 M	<input type="checkbox"/>					
------	--------------------	--------------------------	--	--	--	--	--

23/  
26

Fig.20

200

DEALING SERVICE				/VAL 1998/04/14			
MARKET AMT PTY				YOUR OFFERS			
				123.45 10 1 0 10:03			
				123.50 10 1 0 10:04			
				123.60 10 1 0 10:05			
123.60 * 10 1				YOUR BIDS			
123.50 * 10 1							
123.45 * 10 1							
				/ 2			
10:05 OFFER 123.60 10 MIO MINI 1 H1DN O ***							
				BUY/SELL			
				TOTAL AMT			
				TOTAL TRS			
				AVE. PRICE			
MENU				1993/04/12 10:06 M			

201

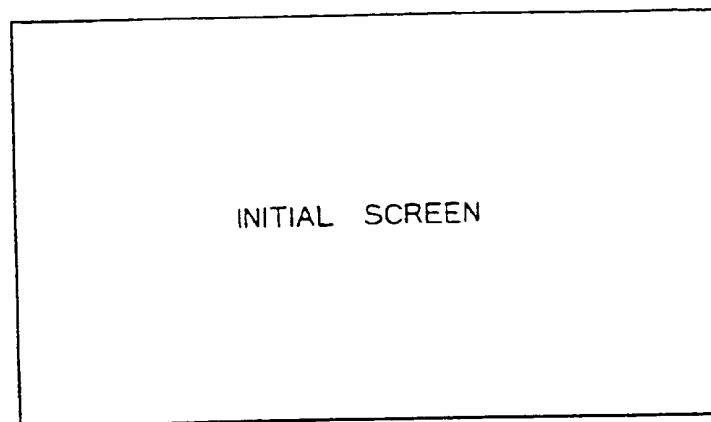


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Fig.21

210

-		DEALING SERVICE		/VAL 1993/04/14	
MARKET AMT PTY		YOUR OFFERS			
		123.45	10	1	0 10:03
		123.50	10	1	0 10:04
		123.60	10	1	0 10:05
123.60					
123.50	-	LEAVE - OUT			
123.45					
		<input type="checkbox"/> CANCEL TIME    10 hour 00 minutes			
		<input type="button" value="EXEC"/> <input type="button" value="QUIT"/>			
10:05 OFFER 123.60 10 MIO MINI 1 HNDN OXXX					
		BUY/SELL TOTAL AMT TOTAL TRS AVE. PRICE			
MENU		1993/04/12 10:06 M			

25/  
26*Fig.22A**Fig.22B*

—	SERVICE LOG-IN	
<input type="checkbox"/>	USER ID	<input type="text"/>
<input type="checkbox"/>	PASSWORD	<input type="text"/>
EXEC		QUIT

26/  
26

Fig.23

— DEALING SERVICE		/VAL 1993/04/14	
MARKET AMT PTY		YOUR OFFERS	231
		123.60 10 1 0 10:05	
		YOUR BIDS	
123.60* 10 1 *			
		/ *	
DONE 10 MIO (1TRS)*** 11:50 SEL 10 MIO 123.50XXXXX DONE 10 MIO (1TRS)*** 11:30 SEL 10 MIO 123.45XXXXX		BUY/SELL	
		TOTAL AMT	
		TOTAL TRS	
		AVE. PRICE	
MENU	1993/04/12 13:15 M		

-1-

DATA PROCESSING SYSTEM

The present invention relates to a data processing system arranged and adapted to provide matching processing between users, in particular to an electronic dealing system used, for example, for  
5 foreign exchange transactions among banks.

At the present time, foreign exchange transactions among banks are performed through the agency of brokers or by direct transactions between the banks. These  
10 transactions are all performed over the telephone (telephone market). Therefore, an electronic dealing system which acts as an computerized electronic agency is desired. In such an electronic dealing system, it is further desired that transactions can be continued  
15 even after the operator of a dealing terminal leaves the terminal.

According to the present invention, there is provided a data processing system arranged and adapted to provide matching processing between users, the  
20 system comprising a computer system arranged to carry out a matching procedure and a plurality of terminals arranged to be coupled to the computer system for the transmission to and from the computer system of user data defining potential matching events, each terminal  
25 including storage means arranged to store user-entered event data, the data processing system including means defining a leave-data function which can be activated by a user from any of said terminals to allow user-entered event data stored in the terminal to continue  
30 to be supplied to the computer after the user has logged out of the computer system, and to allow a matching procedure of said user-entered event data also after the user has logged out.

In a preferred embodiment the system is operable  
35 as a dealing system wherein said user-entered event data is transaction data comprising terms of sale and

terms of purchase, said matching processing being  
transaction processing to match the terms of sale with  
the terms of purchase, and said leave-data function  
being a leave-order function allowing the transaction  
5 matching of said user-entered transaction data to  
continue after the user has logged out.

Thus, in an electronic dealing system embodying  
the invention, the leave-order function enables a  
dealing terminal to continue to place orders on the  
10 market and automatically perform transactions even  
after log-out processing, and thereby\_\_\_\_\_

enables transactions to be safely performed even when the operator is not at the dealing terminal, for example, when the operator has gone home.

Reference is made, by way of example, to the  
5 accompanying drawings in which:-

Fig. 1 is a conceptual view of an example of the  
10 constitution of a foreign exchange transaction in an electronic dealing system to which the present invention is applied;

Fig. 2 is a schematic view of an example of the  
15 overall configuration of an electronic dealing system to which the present invention is applied (part 1);

Fig. 3 is a schematic view of an example of the  
overall configuration of an electronic dealing system to which the present invention is applied (part 2);

Fig. 4 is a view for explaining the processing in  
20 the market by the electronic dealing system to which the present invention is applied;

Fig. 5 is a view for explaining the state of the  
virtual terminal processes (VTP's) in an electronic  
dealing system based on the present invention;

Fig. 6 is a view of an example of a sequence in the  
25 case of a leave-order state in the electronic dealing system of the present invention (part 1);

Fig. 7 is a view of an example of a sequence in the  
30 case of a leave-order state in the electronic dealing system of the present invention (part 2);

Figs. 8A and 8B are views of an example of a control  
sequence in the case of establishment of a transaction in  
the leave-order state in the electronic dealing system of  
the present invention (part 1);

Fig. 9 is a view of an example of a control sequence  
35 in the case of establishment of a transaction in the leave-order state in the electronic dealing system of the

present invention (part 2);

Figs. 10A and 10B are views of an example of a control sequence in processing for canceling a leave-order state at a set time in the electronic dealing system of the present invention (part 1);

Fig. 11 is a view of an example of a control sequence in processing for canceling a leave-order state at a set time in the electronic dealing system of the present invention (part 2);

Figs. 12A and 12B are views of an example of a control sequence in processing for compulsorily resetting a virtual terminal process VTP in the leave-order state in the electronic dealing system of the present invention (part 1);

Fig. 13 is a view of an example of a control sequence in processing for compulsorily resetting a virtual terminal process VTP in the leave-order state in the electronic dealing system of the present invention (part 2);

Fig. 14 is a view of an example of a control sequence in processing for compulsorily resetting a virtual terminal process VTP in the leave-order state in the electronic dealing system of the present invention (part 3);

Fig. 15 is a schematic view of an example of a sequence in processing for compulsorily resetting a virtual terminal process VTP in the leave-order state in the electronic dealing system of the present invention;

Fig. 16 is a flow chart of an example of processing for compulsorily resetting a virtual terminal process VTP in the leave-order state in the electronic dealing system of the present invention;

Fig. 17 is a view of an example of processing for dealing with congestion in a file holding the results of the transactions in the electronic dealing system of the present invention;

Figs. 18A and 18B are views of examples of screens

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displayed on a terminal in the electronic dealing system of the present invention (part 1);

5 Figs. 19A and 19B are views of examples of screens displayed on a terminal in the electronic dealing system of the present invention (part 2);

Fig. 20 is a view of an example of a screen displayed on a terminal in the electronic dealing system of the present invention (part 3);

10 Fig. 21 is a view of an example of a screen displayed on a terminal in the electronic dealing system of the present invention (part 4);

Figs. 22A and 22B are views of examples of screens displayed on a terminal in the electronic dealing system of the present invention (part 5); and

15 Fig. 23 is a view of an example of a screen displayed on a terminal in the electronic dealing system of the present invention (part 6).

#### DESCRIPTION OF THE PREFERRED EMBODIMENTS

20 Below, an embodiment of the electronic dealing system of the present invention will be described with reference to the drawings.

25 Figure 1 is a conceptual view of an example of the constitution of a foreign exchange transaction in an electronic dealing system to which the present invention is applied. In the figure, reference numeral 11 is an electronic dealing system, 12A and 12B are banks, 13 is a central bank, and 14 is a broker (agency).

30 As shown in Fig. 1, a foreign exchange transaction is either performed by a plurality of banks 12A and 12B and a central bank through a broker 14 or else is performed directly between the banks (direct dealing). The banks 12A and 12B are, for example, comprised of interbank dealers, customer dealers, and back offices. The electronic dealing system of the present invention  
35 relates to a foreign exchange transaction (10) performed through a broker 14.

Figure 2 and Fig. 3 are schematic views of an



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example of the overall configuration of an electronic dealing system to which the present invention is applied. In Fig. 2 and Fig. 3, reference numeral 21 is a main center, 22 is a sub-center, 23a and 23b are customer offices, 24 is a dealing part, 25 is a maintenance and operation part, and 26 is an encipher apparatus. Further, reference MAT-H is a matching host serving as the main frame for brokering exchange transactions among customers, SAT-H is a charging and statistic facility for issuing bills for service and managing statistical information etc. in the electronic dealing system, NAM-ST is a general supervisory facility for centrally managing the state of operation of the equipment, and CTL is a subscriber control apparatus set at the subscriber's location for accommodating the lines with the host and controlling the dealing terminals.

The subscriber control apparatus CTL is provided with a management distribution server MDS for controlling the transmission and reception of data between a host and terminals and a video terminal controller VTC for supplying video signals to terminal screens and connecting with existing video terminals. The customer offices 23a and 23b are provided with Confirmation-sheet Automatic Transfer Terminal CAT-T's for storing and printing confirmation sheets (contracts) and a plurality of stand-alone type dealing terminals SAL. Here, the data is enciphered by the data encipher apparatuses 26 so as to maintain confidentiality. Further, the data is divided in time and multiplexed by the time division multiplexer units (TDM). The customer office 23a corresponds, for example, to the dealing room of a domestic bank directly connected to a main center 21, while the customer office 23b corresponds, for example, to the dealing room of a foreign bank connected to the main center 21 and a sub-center 22 provided overseas through an international communication line (satellite communication line, submarine communication

7

cable, etc.)

Figure 4 is a view for explaining the processing in the market by the electronic dealing system to which the present invention is applied.

5 As shown in Fig. 4, first, when an operator places an order (41) through a dealing terminal 40a, that information becomes rate information. Then, for example, if there is a hit (42) from another dealing terminal 40b, the system compares the terms of the transaction (43),  
10 then compares the terms of the correspondent agreements (44), compares the credit lines (45), and outputs the results of the transaction (46) to the dealing terminals 40a and 40b. Here, for example, the operator can set the terms of the transaction when he places the order on the  
15 selling market. The operator determines the selling rate from the orders placed on the selling market. The system outputs changes to all terminals as the rate information (48). The system outputs confirmation sheets (47) through the CAT-T's provided at the customer offices (23a and  
20 23b).

The above description gives a general outline of an electronic dealing system to which the present invention is applied. The features of the electronic dealing system of the present invention will be described in more detail  
25 below.

Figure 5 is a view for explaining the state of the virtual terminal processes VTP's in an electronic dealing system based on the present invention. Here, the "leave order" characterizing the present invention means an  
30 order which a dealing terminal places on the market after the operator logs out.

In Fig. 5, first, in the usual pattern, when an operator starts log-in processing (501) and deal-in processing (502), he becomes able to deal in a  
35 predetermined market. This dealing is stopped when the operator starts deal-out processing (503). He then starts log-out processing (504) so as to end all operations.

That is, in the usual pattern, dealing is possible from the deal-in processing (502) to the deal-out processing (503). Here, when the operator starts the deal-out processing, usually he starts the deal-out processing when there are no orders left. Conversely, in the leave-order pattern of the present invention, explained below, he starts the deal-out processing when there are still orders to be dealt in.

In the leave-order pattern, like with the above usual pattern, when an operator starts log-in processing (51) and deal-in processing (52), he becomes able to deal in a predetermined market. If the operator starts deal-out processing (53) and log-out processing (54) after having set the leave-order function when still able to deal in the market, the dealing terminal continues to place on the market any orders which still exist. Accordingly, when the operator logs out after having set the leave-order function (54), the system establishes transactions automatically for any orders placed on the market which meet the terms of transaction.

In this leave-order pattern, further, if the operator starts the log-out processing (55) from the logged out state with the leave-order function set, the system automatically executes deal-in processing (55) and displays the results of the transactions of the orders which had been placed as leave orders. Like with the usual pattern, further, if the operator starts the deal-out processing (57) without setting the leave-order function, the system stops the dealing and then the operator starts log-out processing (58) to end all operations. That is, in the leave-order pattern, dealing is possible from deal-in processing (52) to deal-out processing (57) when the leave-order function has not been set. When the leave-order function has been set, even deal-out processing (53) and log-out processing (54) are performed, dealing continues. To stop the dealing, the operator has to have started the deal-out processing

(57) without having set the leave-order function.

In this way, according to the electronic dealing system of the present invention, the system can safely continue transactions even when the operator is no longer at the dealing terminal. Accordingly, even when the operator using a certain dealing terminal is not present, the operator can continue to have his orders placed on the market.

In the above description, when the operator starts deal-out processing after having set the leave-order function, the association between the dealing terminal and the subscriber control apparatus (specifically the management distributor server MDS in the subscriber control apparatus CTL) (set by logging in and cut by logging out) is cut, but the individual virtual terminal process VTP continues to be supplied with the dealing service without stopping. Note that the conditions for placement of leave orders may be set as follows: (1) the leave-order function may be set for dealing in a single market, (2) when the system receives a deal-out request with designation of the leave-order function in a state where there are no orders placed, it issues a deal-out response (NG: no orders placed) and awaits a normal deal-out request, and (3) the system has a compulsory virtual terminal process VTP reset function by which an operator can request the compulsory release of the virtual terminal process in the leave-order state, the management distribution server MDS cancels all leave orders of a user when receiving the request for compulsorily resetting the virtual terminal process VTP, and the system stores the content of the file holding the results of transactions in a spill-out message file as a spill-out message.

Figure 6 and Fig. 7 are views of an example of a sequence in the case of a leave-order state in the electronic dealing system of the present invention.

First, as shown in Fig. 6, in the case where the

10

operator has previously set the leave-order function, if there is then a log-in request (601) and deal-in request (603) from the dealing terminal SAL, the management distribution server MDS sends back a log-in acceptance (602) and deal-in response (604) to the dealing terminal to enable dealing. When the operator places an order (605, 613) from the dealing terminal, the system transfers the order through the management distribution server MDS to the matching host MAT, and the matching host MAT sends back a command receipt response (614 and 606) through the management distribution server MDS to the dealing terminal. Here, when the operator normally places an order on the market, the matching host MAT notifies the results of the transaction (placement of order 615, 607) through the management distribution server MDS to the dealing terminal, and the dealing terminal and the management distribution server MDS send back command receipt responses (608, 616) to the management distribution server MDS and the matching host MAT, respectively.

Next, if there is a deal-out request (609) and log-out request (611) from the dealing terminal, the management distribution server MDS sends back a deal-out response (610) and log-out response (612) to the dealing terminal. Here, in the present invention, if there is a deal-out request (609) in the state where there are orders present, the system automatically sets the leave-order function (sets the leave-order function and then logs out the operator). Then, as explained referring to Fig. 5, since the leave-order function is set, the system automatically establishes transactions for orders placed on the market which match the terms of transaction.

Further, as shown in Fig. 7, if there are then a log-in request (617) and deal-in request (619) from the dealing terminal, the management distribution server MDS sends back a log-in acceptance (618) and deal-in response (620) to the dealing terminal to enable dealing. At this

time, the system automatically designates the market for which the leave-order function had been set by the deal-in request (619) (application ID for which leave-order function is set) and displays a screen corresponding to the screen just before that on the dealing terminal. At this time, the dealing terminal simultaneously displays the results of transactions made during the time the leave-order function was set. Note that the above-mentioned application ID is the ID for designating one market from among the various markets.

Figures 8A and 8B and Fig. 9 are views of an example of a control sequence in the case of establishment of a transaction in the leave-order state in the electronic dealing system of the present invention.

First, as shown in Figs. 8A and 8B, if an operator places an order (801, 813, 821) from a dealing terminal VTC in the deal-in state, the system transfers the order through the virtual terminal process VTP and the rate type service (MDS) to the matching host MAT, which then sends back a command receipt response (822, 814, 802). The rate type service (MDS) records the order information (820). Further, when an operator places an order normally on the market, the matching host MAT notifies the results of the transaction (823, 815, 803) to the dealing terminal VTC and the dealing terminal VTC sends back a command receipt response (804, 824) to the rate type service MDS. The rate type service MDS then sends back a command receipt response (824) to the matching host MAT. Here, the matching host MAT notifies the results of transactions for those orders which have been placed (823, 815, 803).

Next, when there is a deal-out request at a dealing terminal VTC (805), the system recognizes if the leave-order function has been set (810). That is, the virtual terminal process VTP requests if there are any pending orders (815) to the rate type service, which refers to the recorded order information (820) and responds if

there are any pending orders (817). At this time, if there is one or more orders present, the system sets the leave-order function and sends back a deal-out response (806) to the dealing terminal VTC. If there is then a log-out request (807) from the dealing terminal VTC, if the virtual terminal process VTP is active (811), the virtual terminal process VTP sends back a log-out response (808) to the dealing terminal VTC which then manages and cuts the terminal connection (809). Here, the system holds the orders placed after the leave-order function has been set in a file in the virtual terminal process VTP for later notification of the results of transactions (818) (sent to the dealing terminal for notification of the results of transactions after the operator has logged in once again). If a transaction is established while the leave-order function is set, that is, if an order placed on the market in the leave-order state meets the terms of transaction and a transaction is established, the matching host MAT notifies the results of the transactions (825, 819) to the virtual terminal process VTP where they are held in a file. Further, at this time, the virtual terminal process VTP sends back a command receipt response (826) through the rate type service to the matching host MAT. Further, the matching host MAT ahead of the rate type service performs processing during the leave-order state in the same way as the usual processing.

Further, as shown in Fig. 9, when there is a log-in request (827) from the dealing terminal VTC in the state with the leave-order function set, the virtual terminal process VTP sends back a log-in response (828). When there is then a deal-in request (829) from the dealing terminal VTC, the virtual terminal process VTP sends back a deal-in response (830). Further, the virtual terminal process VTP takes out results of transactions during the leave-order state (833) from its file holding the results of transactions and notifies them to the dealing terminal

VTC (831). The virtual terminal process VTP receives a command receipt response (832) from the dealing terminal VTC, then recognizes the release of the leave-order function (834). The processing for notification of the results of transactions (831) and the command receipt response (832) is repeated until all the transactions established during the leave-order state finish being sent to the terminal.

As mentioned above, the system is set up so that when it receives a log-in request (827) from a dealing terminal VTC, it automatically displays a screen corresponding to the screen at the time the operator dealt out (805) just before along with the results of the transaction during the leave-order state. That is, in the leave-order state (time when the leave-order function is set), it is possible to display the results of established transactions, orders which have not been filled, and other various types of events.

Figures 10A and 10B and Fig. 11 are views of an example of a control sequence in processing for canceling the leave-order state at a set time in the electronic dealing system of the present invention. The control sequence shown in Figs. 10A and 10B and Fig. 11 basically is the same as that shown in Figs. 8A and 8B and Fig. 9. An explanation will be made only of the portions relating to the processing for cancelation of the leave-order state at a set time.

As shown in Figs. 10A and 10B, when there are one or more orders present in the deal-in state, if there is a deal-out request (901; corresponding to deal-out request (805) in Fig. 8A), the system sets the leave-order function. At this time, the operator inputs the time for cancelation of the leave-order state from the dealing terminal VTC to set this in the virtual terminal process VTP. When the virtual terminal process VTP recognizes that the cancelation time has arrived (902), it sends a request for cancelation of orders (904) to the rate type



service and a request for cancelation of transactions (908) by all cancelation for user (906) to the matching host MAT. Further, the matching host MAT sends back the results of the cancelation of transactions (909) to the  
 5 rate type service, which then deletes the order information (907) and notifies the virtual terminal process VTP of the cancelation of transactions (905). The virtual terminal process (VTP) stores in its file for holding the results of transactions the results of  
 10 transactions in the leave-order state and the processing for cancelation of the leave-order state due to the arrival of the cancelation time (notifies cancelation of transactions).

As shown in Fig. 11, if there is then a log-in request (910) corresponding to log-in request 827 in Fig. 9) from a dealing terminal VTC in the state where the leave-order function is set, the virtual terminal process VTP notifies the dealing terminal VTC of the results of the transactions (911) and of the cancelation of  
 20 transactions (912) kept in its file for holding the results of transactions.

As mentioned above, the system may be constituted not only so that the leave-order function is maintained after being set until the next log-in processing, but  
 25 also so that it is canceled and the orders placed on the market as leave orders are withdrawn when a preset time arrives.

Figures 12A and 12B, Fig. 13, and Fig. 14 are views of an example of a control sequence in processing for  
 30 compulsorily resetting a virtual terminal process VTP in a leave-order state in the electronic dealing system of the present invention. They show the sequence by which another dealing terminal can cancel a leave-order function. The control sequence shown in Figs. 12A and 12B  
 35 to Fig. 14 basically is the same as that shown in Figs. 8A and 8B and Fig. 9. In particular, Figs. 12A and 12B correspond to Figs. 8A and 8B. An explanation will be

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made only of the portions relating to the processing for compulsorily resetting the virtual terminal process VTP in the leave-order state.

5 First, the "processing for compulsorily resetting the virtual terminal process VTP" functions to compulsorily and the processing of the active virtual terminal process VTP of a management distribution server MDS with no Minex In-house Protocol (MIP) association with the terminal at the time when the leave-order  
10 function has been set. Further, as the Protocol Data Unit (PDU), use is made of a compulsory log-out request PDU. The difference between the function for compulsorily resetting the virtual terminal process VTP and a compulsory log-out is the state of whether the user  
15 (operator) is logging in or not and the difference in the method of notification of the cancelation of orders. Here, if an individual virtual terminal process VTP receives a request for resetting, as shown by reference numeral 139 (bottom left in Fig. 12B), it request  
20 cancelation of transactions (135) by ordering a log-out (133), notifies the virtual terminal process VTP of the results of cancelation of transactions (135), then copies the order information to the spill-out message bulletin board of the user.

25 As shown in Fig. 13, in the state where the leave-order function has been set by a certain dealing terminal VTC, for example, when a need arises for another dealing terminal VTC to compulsorily cancel (compulsorily reset the virtual terminal process VTP) the leave-order state  
30 of that dealing terminal VTC (whose operator is not present) due to a sudden change in the rate etc., the other dealing terminal requests compulsory log-out (131) to the requested individual virtual terminal process VTP and orders log-out (133) to that requested individual  
35 virtual terminal process VTP. Further, the requested individual virtual terminal process VTP requests cancelation of transactions (135) to the rate type

service, which requests cancelation of transactions (137) to the matching host MAT. Further, the matching host MAT sends back the results of cancelation of transactions (138) to the rate type service, which then deletes the order information and notifies the virtual terminal process of the cancelation of transactions (136). The virtual terminal process (VTP) keeps the results of the transactions and the processing for another terminal to compulsorily cancel the leave-order state (notify cancelation of transactions) in its file for holding results of transactions during the leave-order state. Further, the virtual terminal process VTP copies the content of the file for holding the results of transactions on to the spill-out message bulletin board (139), then enters the initial state.

Further, as shown in Fig. 14, after the virtual terminal process VTP is compulsorily reset, if there is a log-in request (141) and a deal-in request (143) from a dealing terminal VTC to the same individual virtual terminal process VTP, that virtual terminal process VTP sends back a log-in response (142) and deal-in response (144), but, at this time, it does not display the directly preceding dealing screen, but newly displays the same type of screen (i.e., Fig. 19A) as when the dealing service is started. Further, when there is a message retrieval request (145) from the dealing terminal VTC to the individual virtual terminal process VTP, the virtual terminal process VTP takes out the content copied on the spill-out message bulletin board by its compulsory resetting (such as the results of transactions established before the compulsory resetting) from the spill-out message bulletin board and sends back a message retrieval response (146). This enables the operator to check the results of the transactions.

In this way, when there is a sudden unforeseen change in the rate, even if the operator is not present at a certain dealing terminal, the operator of another

dealing terminal (for example, an adjoining one) can compulsorily reset the virtual terminal process VTP so as to cancel the orders placed on the market by that certain terminal by the leave-order function. Further, it is possible to construct the system so that only a user with a higher ID (for example, an operator with a higher rank) can compulsorily reset the virtual terminal process VTP. In general, the order of rank, from the top down, is a manager, chief dealer, and then dealer.

Figure 15 is a schematic view of an example of a sequence in processing for compulsorily resetting a virtual terminal process VTP in a leave-order state in the electronic dealing system of the present invention.

In the sequence of the PDU at the time of a compulsory reset, when a terminal A sends a compulsory log-out request (151) to the management distribution server MDS, the management distribution server MDS sends back a compulsory log-out response (152). Further, when the terminal B is an active virtual terminal process VTP, the management distribution server MDS requests cancelation of transactions (153) to the matching host MAT, which in turn notifies the management distribution server MDS of the results of cancelation of transactions (154) and releases the active virtual terminal process VTP (155).

Here, the conditions enabling compulsory resetting of a virtual terminal process VTP are (1) that a compulsory virtual terminal process VTP reset function be allowed for the user and (2) that the designated virtual terminal process be an active virtual terminal process VTP.

Figure 16 is a flow chart of an example of processing for compulsorily resetting a virtual terminal process VTP in a leave-order state in the electronic dealing system of the present invention.

As shown in Fig. 16, first, at step 161, it is judged if the rank of the user of the requesting terminal

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allows compulsory log-out service (compulsory resetting of the virtual terminal process VTP) or not. If it allows it, then the routine proceeds to step 162. At step 162, it is judged if there is a user present. If present, the routine proceeds to step 163. At step 163, it is judged if the designated individual virtual terminal process VTP is an active virtual terminal process VTP. When the designated individual virtual terminal process VTP is an active virtual terminal process VTP, the routine proceeds to step 164, where processing is performed for canceling the orders (compulsory resetting of the virtual terminal process VTP) and the routine proceeds to step 165. Here, if the answer is negative (NG) at steps 161 to 163, the routine proceeds to step 168, where a compulsory log-out response (NG) is sent out.

At step 165, the results of the cancelation and the results of the transactions are saved on the spill-out message bulletin board, then the routine proceeds to step 166. Here, the information saved on the spill-out message bulletin board at step 165 (results of cancelation and results of transactions) can be read out at the request of the user (operator). At step 166, the log of the compulsory log-out is acquired and then the routine proceeds to step 167, where a compulsory log-out response (OK) is sent out.

Here, the information on the establishment of transactions and the information on cancelation of orders etc. at the active virtual terminal process VTP are kept (stored) in the file for holding the results of transactions, but the user (operator) can determine changes in the orders when logging out from the stored information sent out at the time of logging in. Further, the PDU stored in the file holding the results of transactions includes (1) notifications of the results of transactions, (2) notifications of cancelation of transactions, and (3) requests for status confirmation.

Figure 17 is a view of an example of processing for

dealing with congestion in a file for holding results of transactions in the electronic dealing system of the present invention.

5 As shown in Fig. 17, the file for holding the results of transactions stores messages until the resources of the system become congested (until the warning level is reached, for example, until 80 percent of capacity is reached). Here, when the system detects the congested state and the warning level is reached, the  
10 system cancels all orders and stores notifications of cancelation of transactions in the file for holding the results of transactions. Further, when a command to cancel all transactions and notifications of the results of transactions and notifications of cancelation of  
15 transactions or the like cross, the notifications of results of transactions and other messages are spilled out.

20 Figures 18A and 18B to Fig. 23 are views of examples of the display screens of a dealing terminal in the electronic dealing system of the present invention.

First, when the operator turns on the power of the dealing terminal, the terminal displays the initial screen shown in Fig. 18A. When he then depresses any key, the terminal displays the log-in screen shown in Fig. 18B  
25 (log-in window screen). When he inputs his user ID, password, etc. in the log-in window screen of Fig. 18B and logs in, that is, when the log-in is accepted, the screen becomes the management screen shown in Fig. 19A. Further, when he initiates log-in processing at the  
30 management screen of Fig. 19A, more particularly, when a deal-in response is received, the screen becomes the dealing service screen shown in Fig. 19B and dealing can be performed.

Next, as shown in Fig. 20, the operator places  
35 orders on the market, for example, places three offers (200), by an offer command (OFFER) on the dealing service screen of Fig. 19B. More specifically, in this



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function had been set, the terminal changes from the screen of Fig. 22A to the log-in window screen of Fig. 22B. When the operator then inputs his user ID and password etc. so as to log in, the terminal does not display the management screen corresponding to Fig. 19A, but automatically performs deal-in processing and displays the dealing service screen shown in Fig. 23. The screen shown in Fig. 23 corresponds to the screen shown in Fig. 20 at the time of the deal-out processing just before. In Fig. 23, however, the results of the transactions established during the leave-order state are displayed. That is, the screen displays the results of transactions (230) established during the leave-order state (time during which the leave-order function is set), more specifically, the information (230) showing that the offer order of the price of "123.45" and the amount of "10" and the offer order of the price of "123.50" and the amount of "10" have been hit and the orders (231) remaining without establishment of transactions, more specifically, the information (231) showing that the offer order of the price of "123.60" and the amount of "10" remains.

In the above explanation, the electronic dealing system of the present invention was explained with reference to foreign exchange transactions, but the invention can be applied to various other types of transactions as well. Further, the display screens are not limited to those shown in Figs. 18A and 18B to Fig. 23 and can be modified in various ways.

As explained in detail above, according to the electronic dealing system of the present invention, since the electronic dealing system is given a "leave-order" function by which a dealing terminal can continue to place orders on the market and automatically perform transactions even after the operator has logged out, transactions can be continued safely even when the operator is not present.



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CLAIMS:

1. A data processing system arranged and adapted to provide matching processing between users, the system comprising a computer system arranged to carry out a matching procedure and a plurality of terminals arranged to be coupled to the computer system for the transmission to and from the computer system of user data defining potential matching events, each terminal including storage means arranged to store user-entered event data, the data processing system including means defining a leave-data function which can be activated by a user from any of said terminals to allow user-entered event data stored in the terminal to continue to be supplied to the computer after the user has logged out of the computer system, and to allow a matching procedure of said user-entered event data also after the user has logged out.
2. A data processing system as set forth in claim 1, the system being operable as a dealing system wherein said user-entered event data is transaction data comprising terms of sale and terms of purchase, said matching processing being transaction processing to match the terms of sale with the terms of purchase, and said leave-data function being a leave-order function allowing the transaction matching of said user-entered transaction data to continue after the user has logged out.
3. A data processing system as set forth in claim 2, wherein the leave-order function is set by deal-out processing in a state where at least one order is present in the dealing terminal.
4. A data processing system as set forth in claim 3, wherein when log-in processing is executed

after deal-out processing is executed with the leave-order function set, the system automatically executes deal-in processing as well and displays a screen corresponding to the screen at the time of the deal-out processing together with the results of transactions in the time when the leave-order function had been set.

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5. A data processing system as set forth in claim 2, 3 or 4, wherein provision is made of a facility for setting the time for cancelation of the leave-order function.

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6. A data processing system as set forth in any of claims 2 to 5, wherein provision is made of a facility for cancelation of the leave-order function by another dealing terminal.

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7. A data processing system as set forth in any of claims 2 to 6, wherein the leave-order mode is entered by the execution of deal-out processing in a state with orders remaining in the dealing system.

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8. A data processing system as set forth in any of claims 2 to 7, which performs log-in processing and deal-in processing to enable dealing in a predetermined market, sets the leave-order function and performs deal-out processing and log-out processing in the state where this dealing can be executed and continues to place any orders which still exist on the market.

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9. A data processing system as set forth in claim 8, which executes log-out processing while setting the leave-order function and automatically establishes transactions for orders placed on the market conditional on the matching of terms of transaction.

10. A data processing system as set forth in claim 8, which executes log-in processing from the log-out state with the leave-order function set so as to automatically perform deal-in processing and displays the results of the transactions of the orders which had been placed as leave orders.

11. A data processing system as set forth in claim 7, which continues the dealing even if the deal-out processing and log-out processing are performed until the deal-out processing is performed without the leave-order function being set.

12. A data processing system as set forth in claim 7, wherein when executing deal-out processing while setting a leave-order function, the association between the dealing terminal and a subscriber control apparatus, which is set by logging in and cut by logging out, is cut, but the individual virtual terminal process continues to supply the dealing service without stopping.

13. A data processing system as set forth in claim 12, which receives a deal-out request and a log-out request from the dealing terminal, returns a deal-out response and a log-out response from said subscriber control apparatus to the dealing terminal and, at that time, carries out the deal-out request in the state where there are orders present so as to automatically set the leave-order function, whereby if the leave-order function is set, transactions are automatically established for orders placed on the market which match the terms of transaction.

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